

FIGURE 4.10 A UML class diagram corresponding to the EER diagram in Figure 4.7, illustrating UML notation for specialization/generalization.

In database design, we are mainly concerned with specifying concrete classes whose collections of objects are permanently (or persistently) stored in the database. The bibliographic notes at the end of this chapter give some references to books that describe complete details of UML. Additional material related to UML is covered in Chapter 12, and object modeling in general is further discussed in Chapter 20.

4.7 RELATIONSHIP TYPES OF DEGREE HIGHER THAN TWO

In Section 3.4.2 we defined the **degree** of a relationship type as the number of participating entity types and called a relationship type of degree two *binary* and a relationship type of degree three *ternary*. In this section, we elaborate on the differences between binary

and higher-degree relationships, when to choose higher-degree or binary relationships, and constraints on higher-degree relationships.

4.7.1 Choosing between Binary and Ternary (or Higher-Degree) Relationships

The ER diagram notation for a ternary relationship type is shown in Figure 4.11a, which displays the schema for the `SUPPLY` relationship type that was displayed at the instance level in Figure 3.10. Recall that the relationship set of `SUPPLY` is a set of relationship instances (s, j, p) , where s is a `SUPPLIER` who is currently supplying a `PART` p to a `PROJECT` j . In general, a relationship type R of degree n will have n edges in an ER diagram, one connecting R to each participating entity type.

Figure 4.11b shows an ER diagram for the three binary relationship types `CAN_SUPPLY`, `USES`, and `SUPPLIES`. In general, a ternary relationship type represents different information than do three binary relationship types. Consider the three binary relationship types `CAN_SUPPLY`, `USES`, and `SUPPLIES`. Suppose that `CAN_SUPPLY`, between `SUPPLIER` and `PART`, includes an instance (s, p) whenever supplier s can supply part p (to any project); `USES`, between `PROJECT` and `PART`, includes an instance (j, p) whenever project j uses part p ; and `SUPPLIES`, between `SUPPLIER` and `PROJECT`, includes an instance (s, j) whenever supplier s supplies some part to project j . The existence of three relationship instances (s, p) , (j, p) , and (s, j) in `CAN_SUPPLY`, `USES`, and `SUPPLIES`, respectively, does not necessarily imply that an instance (s, j, p) exists in the ternary relationship `SUPPLY`, because the meaning is different. It is often tricky to decide whether a particular relationship should be represented as a relationship type of degree n or should be broken down into several relationship types of smaller degrees. The designer must base this decision on the semantics or meaning of the particular situation being represented. The typical solution is to include the ternary relationship plus one or more of the binary relationships, if they represent different meanings and if all are needed by the application.

Some database design tools are based on variations of the ER model that permit only binary relationships. In this case, a ternary relationship such as `SUPPLY` must be represented as a weak entity type, with no partial key and with three identifying relationships. The three participating entity types `SUPPLIER`, `PART`, and `PROJECT` are together the owner entity types (see Figure 4.11c). Hence, an entity in the weak entity type `SUPPLY` of Figure 4.11c is identified by the combination of its three owner entities from `SUPPLIER`, `PART`, and `PROJECT`.

Another example is shown in Figure 4.12. The ternary relationship type `OFFERS` represents information on instructors offering courses during particular semesters; hence it includes a relationship instance (i, s, c) whenever `INSTRUCTOR` i offers `COURSE` c during `SEMESTER` s . The three binary relationship types shown in Figure 4.12 have the following meanings: `CAN_TEACH` relates a course to the instructors who can teach that course, `TAUGHT_DURING` relates a semester to the instructors who taught some course during that semester, and `OFFERED_DURING` relates a semester to the courses offered during that semester by any instructor. These ternary and binary relationships represent different information, but certain constraints should hold among the relationships. For example, a relationship instance (i, s, c) should not exist in `OFFERS` unless an instance (i, s) exists in `TAUGHT_DURING`,

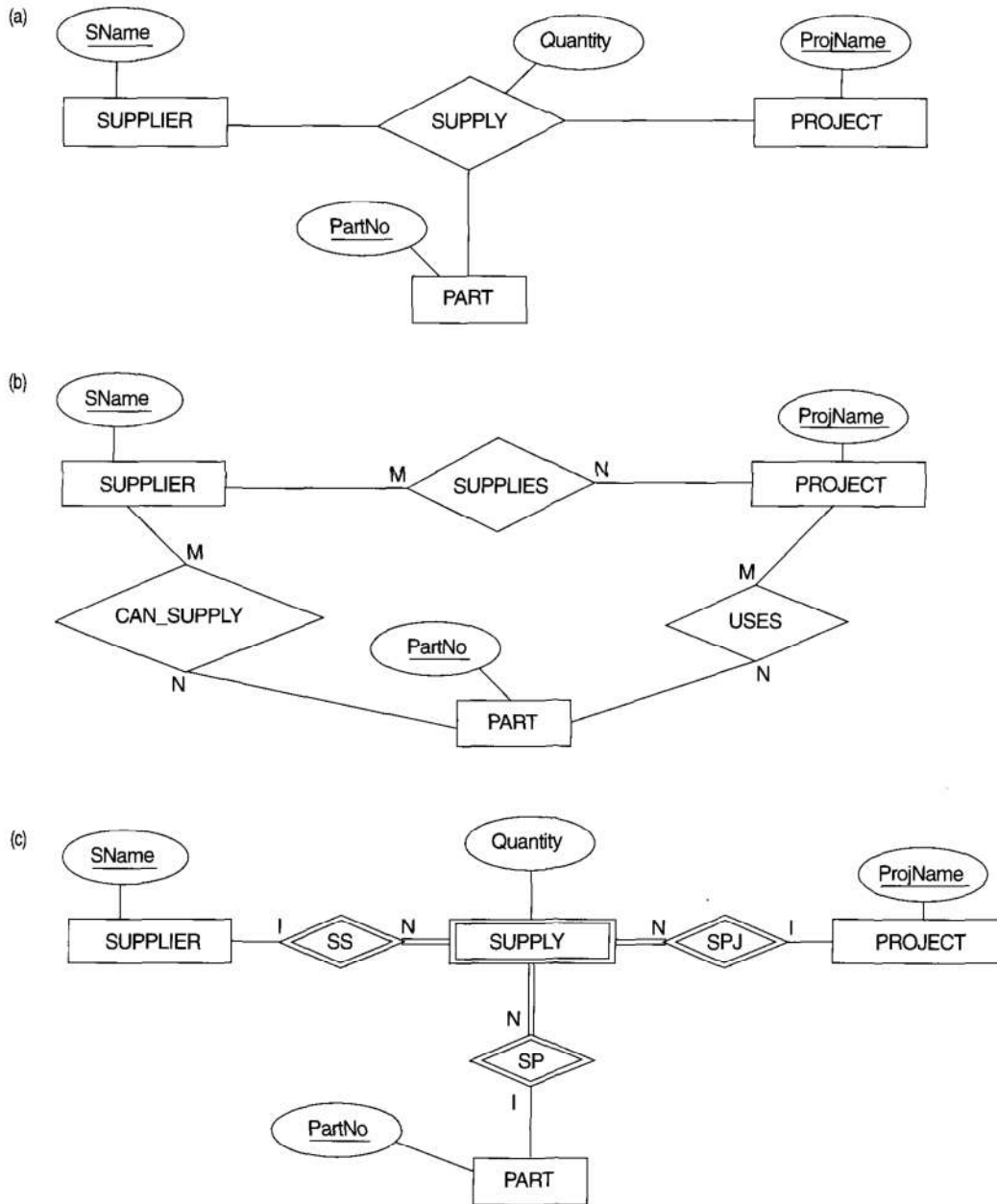


FIGURE 4.11 Ternary relationship types. (a) The SUPPLY relationship. (b) Three binary relationships not equivalent to SUPPLY. (c) SUPPLY represented as a weak entity type.

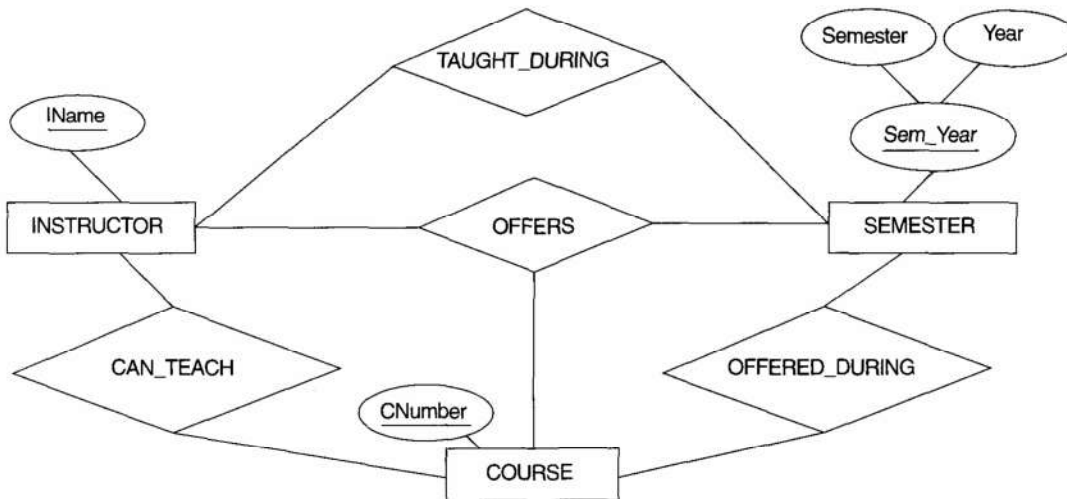


FIGURE 4.12 Another example of ternary versus binary relationship types.

an instance (s, c) exists in OFFERED_DURING, and an instance (i, c) exists in CAN_TEACH. However, the reverse is not always true; we may have instances (i, s) , (s, c) , and (i, c) in the three binary relationship types with no corresponding instance (i, s, c) in OFFERS. Note that in this example, based on the meanings of the relationships, we can infer the instances of TAUGHT_DURING and OFFERED_DURING from the instances in OFFERS, but we cannot infer the instances of CAN_TEACH; therefore, TAUGHT_DURING and OFFERED_DURING are redundant and can be left out.

Although in general three binary relationships cannot replace a ternary relationship, they may do so under certain *additional constraints*. In our example, if the CAN_TEACH relationship is 1:1 (an instructor can teach one course, and a course can be taught by only one instructor), then the ternary relationship OFFERS can be left out because it can be inferred from the three binary relationships CAN_TEACH, TAUGHT_DURING, and OFFERED_DURING. The schema designer must analyze the meaning of each specific situation to decide which of the binary and ternary relationship types are needed.

Notice that it is possible to have a weak entity type with a ternary (or n -ary) identifying relationship type. In this case, the weak entity type can have *several* owner entity types. An example is shown in Figure 4.13.

4.7.2 Constraints on Ternary (or Higher-Degree) Relationships

There are two notations for specifying structural constraints on n -ary relationships, and they specify different constraints. They should thus *both be used* if it is important to fully specify the structural constraints on a ternary or higher-degree relationship. The first

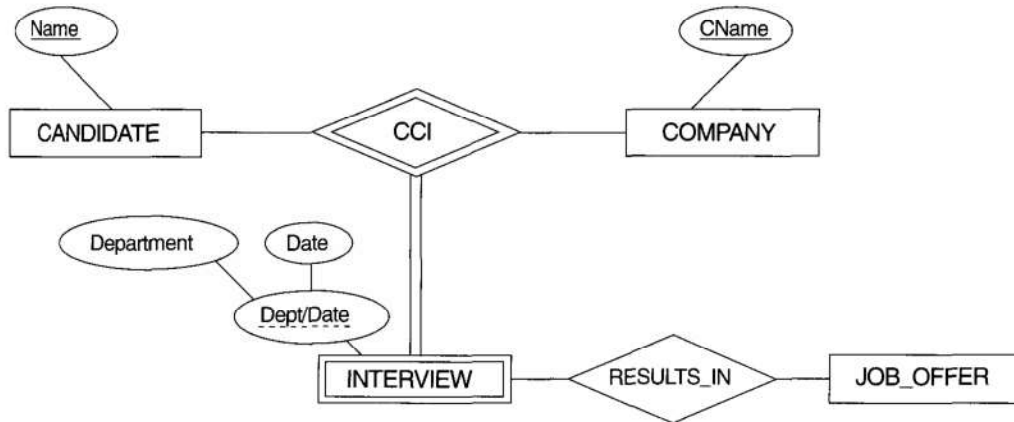


FIGURE 4.13 A weak entity type *INTERVIEW* with a ternary identifying relationship type.

notation is based on the cardinality ratio notation of binary relationships displayed in Figure 3.2. Here, a 1, M, or N is specified on each participation arc (both M and N symbols stand for *many* or *any number*).¹² Let us illustrate this constraint using the *SUPPLY* relationship in Figure 4.11.

Recall that the relationship set of *SUPPLY* is a set of relationship instances (s, j, p) , where s is a *SUPPLIER*, j is a *PROJECT*, and p is a *PART*. Suppose that the constraint exists that for a particular project-part combination, only one supplier will be used (only one supplier supplies a particular part to a particular project). In this case, we place 1 on the *SUPPLIER* participation, and M, N on the *PROJECT*, *PART* participations in Figure 4.11. This specifies the constraint that a particular (j, p) combination can appear *at most once* in the relationship set because each such (project, part) combination uniquely determines a single supplier. Hence, any relationship instance (s, j, p) is *uniquely identified* in the relationship set by its (j, p) combination, which makes (j, p) a key for the relationship set. In this notation, the participations that have a one specified on them are not required to be part of the identifying key for the relationship set.¹³

The second notation is based on the (\min, \max) notation displayed in Figure 3.15 for binary relationships. A (\min, \max) on a participation here specifies that each entity is related to at least \min and at most \max relationship instances in the relationship set. These constraints have no bearing on determining the key of an n -ary relationship, where $n > 2$,¹⁴ but specify a different type of constraint that places restrictions on how many relationship instances each entity can participate in.

12. This notation allows us to determine the key of the *relationship relation*, as we discuss in Chapter 7.

13. This is also true for cardinality ratios of binary relationships.

14. The (\min, \max) constraints can determine the keys for binary relationships, though.